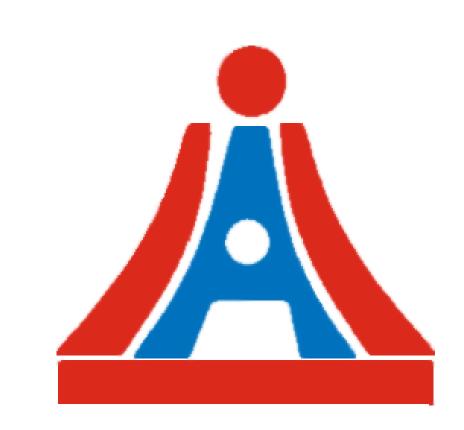


## Anytime Heuristic for Weighted Matching Through Altruism-Inspired Behavior

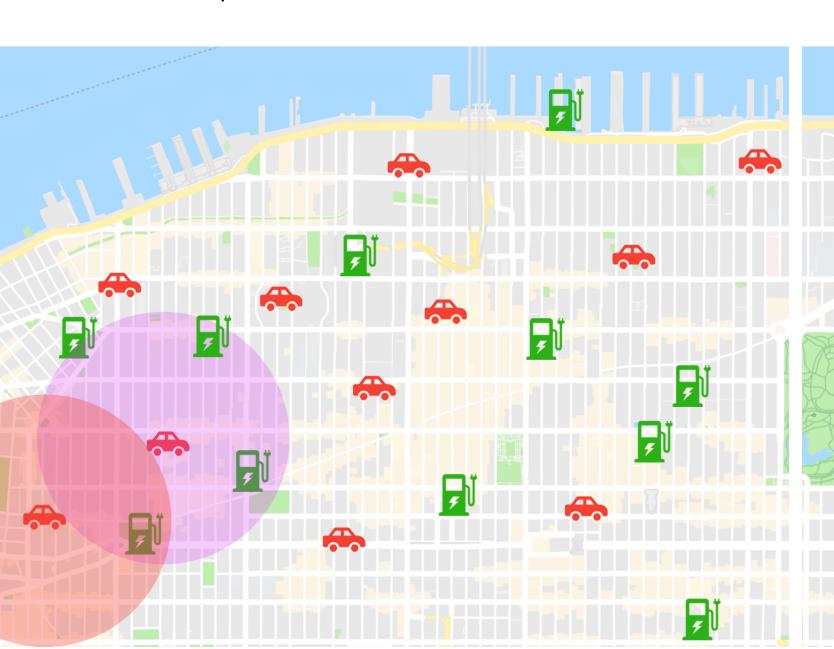
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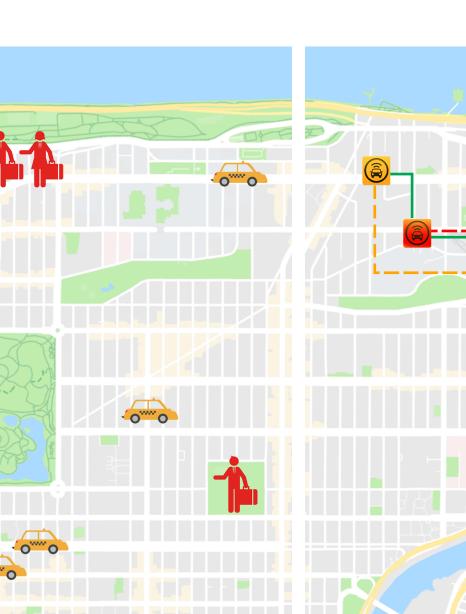


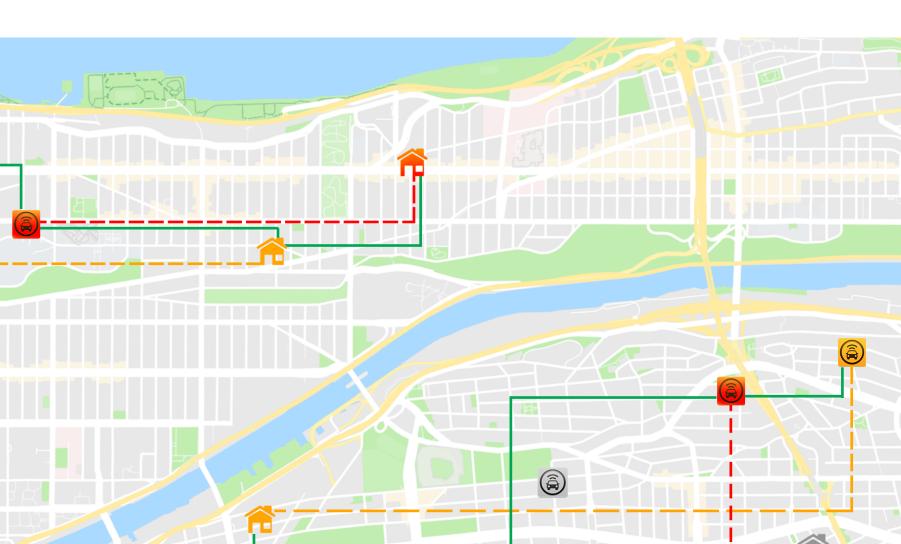
#### Motivation

Charging / Parking Spaces:



Taxi-Passenger Matching:





Ride-Sharing:

#### The Assignment Problem:

- Weighted bipartite graph
- N agents compete for R resources
- Each agent is interested in a subset  $(\mathcal{R}^n)$  of the resources
- Goal: maximize the social welfare (sum of utilities)

#### Applications:

- Role allocation (e.g., robot team formation)
- Task assignment (e.g., taxi-passenger matching)
- Resource allocation (e.g., parking/charging spaces)

#### Challenges in Real World:

Simulation Results

0.2

0.85

Percentage of the trip time each request waits (q)

0.78

Table 1: Empirical Competitive Ratio of ALMA

 $(0, \infty, 0.1)$   $(0, \infty, 0.5)$   $(0, \infty, 1.0)$ 

0.69

0.67

0.1

0.79

- Distributed nature
- Information restrictive (lack of communication/responsiveness, partial feedback)
- Large scale

Hundreds of thousands of autonomous agents (e.g., IoT devices, intelligent infrastructure, autonomous vehicles, etc.)

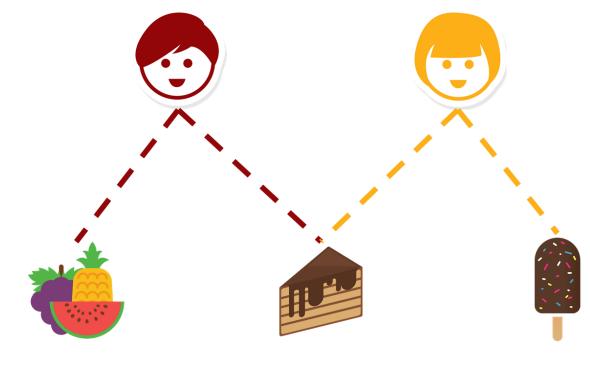
Existing algorithms require: (i) runtime that increases with the total problem size, even if the agents are interested in a few resources, (ii) significant amount of inter-agent communication.

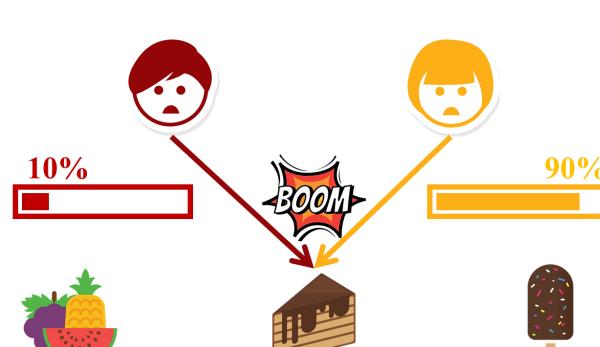
Need for fast convergence to allocations of high social welfare.

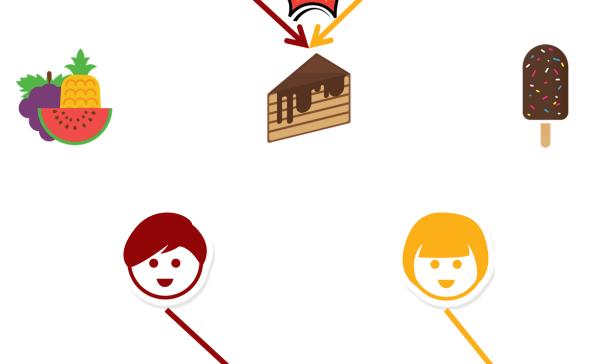
Humans are routinely called upon to coordinate in large scale, and under dynamic and unpredictable demand. Driving factor: principle of altruism.

#### ALMA: ALtruistic MAtching heuristic

Agents make decisions locally, based on (i) the contest for resources that *they* are interested in, (ii) the agents that are interested in the *same* resources. If each agent is interested in only a *subset* of the total resources, ALMA converges in *constant* time. The same is not true for other algorithms, which require time polynomial in the *total* number of agents/resources, even if the aforementioned condition holds. The condition holds by default in many real-world applications; agents have only local knowledge of the world, there is typically a cost associated with acquiring a resource, or agents are simply only interested in resources in their vicinity (e.g., urban environments).







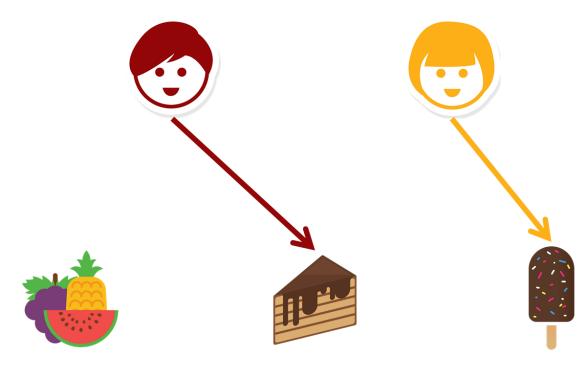


- Go over the set of preferred resources  $\mathcal{R}^n \subseteq \mathcal{R}$  sequentially.
- If collision, back-off with probability that depends on the utility loss of switching to the remaining resources (e.g., P(loss) = 1 loss).

$$loss_n^i = rac{\sum\limits_{j=i+1}^k u_n(r_i) - u_n(r_j)}{k-i}$$

 $\begin{cases} \text{Good alternatives} \to \text{More likely to back-off} \\ \text{No good alternatives} \to \text{Less likely to back-off} \end{cases}$ 

If back-off  $\rightarrow$  select an alternative resource and examine its availability.



#### Altruism-Inspired Behavior Give up a resource:

- To someone who values it more, to increase the social welfare
- To be nice to others; especially when there are equally good alternatives
- → Faster convergence outweighs the loss in utility.

**Theorem 1** (Convergence Speed).

$$\mathcal{O}\left(\max_{n'\in\cup_{r\in\mathcal{R}^n}\mathcal{N}^r} R^{n'} \frac{2-p_n^*}{2(1-p_n^*)} \left(\frac{1}{p_n^*} \log(\max_{r\in\mathcal{R}^n} N^r) + \max_{n'\in\cup_{r\in\mathcal{R}^n}\mathcal{N}^r} R^{n'}\right)\right) = constant$$

#### Highlights:

- i) Decentralized, completely uncoupled, no communication, only partial feedback
- ii) Constant in the total problem size convergence time, under reasonable assumptions on the preference domain
- iii) High social welfare in a variety of scenarios: synthetic and real data, time constraints, on-line settings

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### Read the Full Paper: